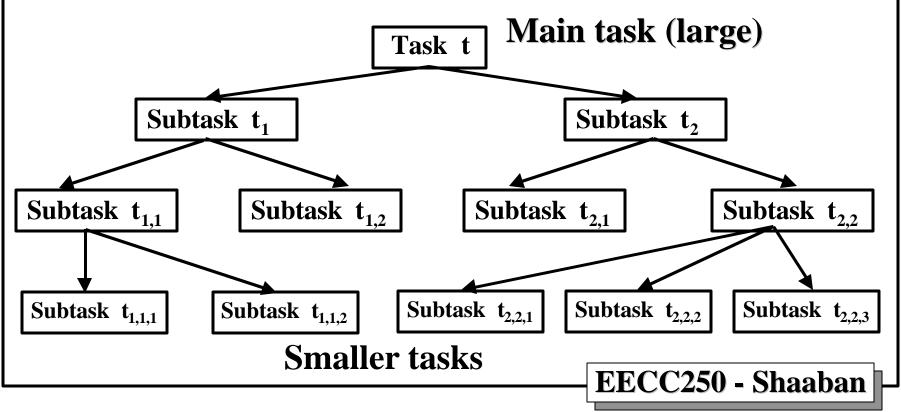
The Software Design Process

The design process of software for microprocessor systems should incorporate the following 5 key concepts:

- Top-Down Design:
- Modular Design
- Testability
- Recoverability
- Structured Programming

Top-down Design

- Programming by step-wise refinement; i.e., decompose a large complex project or task into smaller, more manageable components or subtasks.
- Iterative process that separates the goals of the program from the methods of achieving them.
- Usually accompanied by bottom-up coding.



System Specification

- Before a system (software or hardware) can be designed, it must be *specified*.
- A system specification provides the statement of the goals that a system should achieve.
- The programmer should always validate the end product against these goals.
- A tightly-specified system covering many possible cases is usually more reliable than a loosely specified system.
- It may also be useful to specify "non-goals"; i.e., things that the system is not required to do.

Modular Design

- A software module is concerned with a single, logically coherent task.
- Modules can be "plugged into the system" and can be supplied by different software vendors.
- The internal operation of the module is not significant; only its inputs and outputs.
- Modules can be tested separately from the main system.
- Coupling indicates how information is shared among modules. Strongly coupled modules share common data which is generally undesirable.
- The strength of a module is related to whether or not it performs a single function. Strong modules are easier to test and replace.

Single Entry Point

Module

Module Code (i.e. program)

Local Data Storage (cannot be accessed outsides the module)

> Single Exit Point

Testability

- Testing is done by examining the state of a system at various key points in its lifetime. This can be implemented through the use of breakpoints.
- Bottom-up testing:
 - Involves testing the lowest-level components of a system first.
 - Starts at the lowest level and keeps moving to higher levels.
 - Complete when the highest level of the system has been tested.
 - Requires writing a test driver for the component to be tested.
- Top-down testing:
 - Involves testing the highest levels first.
 - Helps to spot major design problems early.
 - Does not require a test driver for components; but instead requires stubs to represent the lower level modules.
- White Box versus Black Box testing:
 - Black Box testing means that the inner workings are totally unknown;
 thus, all possible inputs and outputs must be tested.
 - White Box testing means that the inner workings are known; this knowledge can be used to limit the amount of testing required.

Recoverability

- Recoverability or exception handling is the ability of a system to cope with erroneous data and to recover from certain classes of errors.
- A poor recovery mechanism may be worse than none at all.

Structured Programming

- Purpose of structure programming:
 - Improve programmer productivity;
 - Make programs easier to read;
 - Yield more reliable programs.
- All programs can be constructed from three fundamental components:
 - Sequence:
 - A linear list of actions that are executed in order.
 - Looping Mechanism:
 - Permits a sequence to be carried out a number of times.
 - Decision Mechanism:
 - Allows one of two courses of action to be taken.

The Conditional Structure

For the purpose of the following, assume that 'L' is a logical condition whose result 'B' is stored in register D0 and S, S1 and S2 are sequences.

*

```
* IF L THEN S
```

TST.B D0 Test the lower-order byte of D0

BEQ ENDIF If not true, then skip the sequence

S

ENDIF ...

* IF L THEN S1 ELSE S2

*

TST.B D0 Test the lower-order byte of D0

BEQ ELSE If not ture, then proceed to the else sequence

S1 Execute the S1 sequence

BRA ENDIF Skip the else statement

ELSE S2 Execute the S2 sequence

ENDIF ...

The CASE Statement

```
*
     CASE I OF
*
                    I1: S1
*
                    I2: S2
                     In: Sn
         MOVE
                  I.D0
                             Move the variable to D0 for testing
         CMP
                 I1.D0
                             Check if it is I1
                 ACT1
         BEO
         CMP
                I2,D0
                             Check if it is I2
         BEQ
                 ACT2
         CMP
                             Check if it is In
                 In,D0
         BEO
                 ACTn
                 ERROR
         BEQ
ACT1
         S1
                             Execute the statement for I1
         BRA
                 ENDCASE
ACT2
         S2
                             Execute the statement for I2
          BRA
                 ENDCASE
ACTn
                             Execute the statement for In
          Sn
                 ENDCASE
          BRA
ERROR
                             Handle a value out of range
ENDCASE
```

The CASE Statement

If the conditions can be converted to a sequence of integer numbers, then the CASE statement is more efficiently handled by a jump table:

	CLR.L DO	Clear all bits of D0
	LEA JUMPTAB,A0	Store the address of the jump table
	MOVE I,D0	Move the variable to D0 for testing
	CMP I1,D0	Check the bottom of the range
	BLO ERROR	Error if less the lowest value
	CMP In,D0	Check the top of the range
	BCS ERROR	Error if higher than highest value
	SUB I1,D0	Get the offset from the first condition
	ASL.L #2,D0	Multiply by 4, addresses are long words
	MOVEA.L (A0,D0),A0	Get the address of the action
	$\mathbf{JMP} \qquad \qquad (\mathbf{A0})$	Jump to the appropriate action
	 D.C.I. A.C.E.I	T
JUMPTAB	DC.L ACT1	First action
	DC.L ACT2	Second action
	DC.L ACTn	N'th action

ACT1	S1	Execute the statement for I1
	BRA ENDCASE	
ACT2	S2	Execute the statement for I2
	BRA ENDCASE	
	•••	
ACTn	Sn	Execute the statement for In
	BRA ENDCASE	
ERROR	Handle a value out of range	
ENDCASE		EECC250 - Shaabar

Looping Mechanisms

```
FOR I = N1 TO N2 DO S
*
       MOVE.B
                 #N1,D0
                              D0 is the loop counter
NEXT
       CMP.B
                 #N2,D0
                              Check if the end of the loop
                              Quit the loop if counter too high
       BHI
                 ENDLOOP
       S
                              Execute the sequence
       ADDO
                 #1,D0
                              Increment the loop counter
       RRA
                 NEXT
ENDLOOP
*
        FOR I = N DOWNTO 0
       MOVE.W
                   #N,D0
                              D0 is the loop counter
                 ENDLOOP
                              Skip loop if less than 0
       BMI
```

Execute the sequence

Decrement D0 and loop back

NEXT

S

DBRA

D0,NEXT

DBcc Test Condition, Decrement, and Branch (M68000 Family)

DBcc

Operation: If Condition False

Then (Dn – 1 \rightarrow Dn; If Dn \neq – 1 Then PC + d_n \rightarrow PC)

Assembler

Syntax: DBcc Dn, < label >

Attributes: Size = (Word)

Description: Controls a loop of instructions. The parameters are a condition code, a data register (counter), and a displacement value. The instruction first tests the condition for termination; if it is true, no operation is performed. If the termination condition is not true, the low-order 16 bits of the counter data register decrement by one. If the result is − 1, execution continues with the next instruction. If the result is not equal to − 1, execution continues at the location indicated by the current value of the program counter plus the sign-extended 16-bit displacement. The value in the program counter is the address of the instruction word of the DBcc instruction plus two. The displacement is a twos complement integer that represents the relative distance in bytes from the current program counter to the destination program counter. Condition code cc specifies one of the following conditional tests (refer to Table 3-19 for more information on these conditional tests):

DBcc

Test Condition, Decrement, and Branch (M68000 Family)

DBcc

Mnemonic	Condition	
CC(HI)	Carry Clear	
CS(LO)	Carry Set	
EQ	Equal	
F	False	
GE	Greater or Equal	
GT	Greater Than	
н	High	
LE	Less or Equal	

Mnemonic	Condition	
LS	Low or Same	
LT	Less Than	
MI	Minus	
NE	Not Equal Plus	
PL		
Т	True	
VC	Overflow Clear	
vs	Overflow Set	

Condition Codes:

Not affected.

DBF Dn,<label>
decrement Dn and branch if Dn has not reached -1
Some assemblers allow DBRA instead of DBF

Looping Mechanisms

*

* WHILE L DO S

*

REPEAT TST.B D0

BEO ENDLOOP

S

BRA REPEAT

Test if the condition still true

If false, then quit

Execute the sequence

Repeat the loop

ENDLOOP

*

* REPEAT S UNTIL L

*

NEXT S

TST.B D0

BNE NEXT

Execute the sequence

Test the value of the condition

If not true, then loop again

ENDLOOP

Pseudocode, or Program Design Language (PDL)

PDL is simply a methodology for expressing the steps of a program before it is translated into assembler. It has the following characteristics:

- A compromise between a high-level language description and assembly language.
- Facilitates the production of reliable code by providing an intermediate step.
- Shares some of the features of high-level languages but without their complexity.
- Provides a shorthand notation for the precise description of algorithms.
- Can be extended to deal with specific tasks.

Example: Comparing two strings

Problem Statement:

A sequence of ASCII characters is stored at memory location \$600 onward (each character one byte). A second string of equal length is stored at memory location \$700 onward. Each string ends with the character \$0D (i.e. carriage return). Write a program to determine if these two strings are equal. If they are identical, then place an \$FF in D0; otherwise, place the value \$00 in D0.

First Level PDL - Indicates what to do:

Match := false

REPEAT

Read a pair of characters

IF they do not match then EXIT

UNTIL a character = \$0D

Match := true

EXIT

Example (continued)

Second Level PDL - **Elaborates on how to do it:**

Match := false

Set pointer1 to point to String1

Set pointer2 to point to String2

REPEAT

Read the character pointed at by String1

Compare with the character pointed at by String2

IF they do not match, THEN EXIT

Pointer1 := Pointer1 + 1

Pointer2 := Pointer2 + 1

UNTIL Character = \$0D

Match := true

EXIT

```
Example: First Assembly Program
               Error Flag
*
       \mathbf{D0}
               Pointer to string 1
       \mathbf{A0}
               Pointer to string 2
       A1
          ORG$400
                                      Start of program
                      #$00,D0
                                      Set the flag to fail
          MOVE.B
                      #$600,A0
                                      A0 points to string1
          MOVEA.L
                                      A1 points to string 2
          MOVEA.L
                      #$700,A1
REPEAT MOVE.B
                       (A0),D1
                                      Get a character from string1
                      (A1),D1
                                      Compare with string2 character
         CMP.B
                                      If characters are different exit
                      EXIT
         BNE
         ADDA.L
                      #1,A0
                                      If the two characters are the
         ADDA.L
                      #1,A1
                                      same point to the next pair
                      #$0D,D1
                                      Test for end of strings
         CMP .B
                                      If not compare next pair
                      REPEAT
         BNE
                      #$FF,D0
                                      ELSE Set flag to success
         MOVE.B
EXIT
         STOP
                      $600
         ORG
Pointer1 DS.B
                       <length of string1>
         ORG
                      $700
Pointer1 DS.B
                      <length of string2>
                                              EECC250 - Shaaban
```

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Example: Refined Assembly Code

-	_		
Car_Ret	EQU	\$0D	
	ORG	\$400	Start of program
	CLR.B	$\mathbf{D0}$	Set the flag to fail
	LEA	Pointer1,A0	A0 points to string1
	LEA	Pointer2,A1	A0 points to string2
REPEAT	MOVE.B	(A0),D1	Get character from string1
	CMP.B	(A1),D1	Compare it with string2
	BNE	EXIT	If different then EXIT
	LEA	1(A0),A0	Point to next pair of characters
	LEA	1(A1),A1	
	CMP.B	#Car_Ret,D1	Test for end of strings
	BNE	REPEAT	If not then compare next pair
	MOVE.B	# \$FF,D 0	ELSE set D0 to success
EXIT	STOP		
	ORG	\$600	
Pointer1	DS.B	<length of="" string1=""></length>	
	ORG	\$700	
Pointer1	DS.B	<le>dength of string2></le>	
		_	_