

Recursive Subroutine Calls Example

The purpose of this example is to examine how all parameters, local variables, return addresses, and frame pointers are stored on the stack when a main program calls a procedure "Process" as well as when the procedure calls itself again in a recursion. We assume the following:

- The stack pointer initially has the value value \$00000F00 just before Process is invoked (before any parameters are pushed onto the stack).
- Array "X", "Y", "Z" and "ALPHA" are passed by reference.
- Parameter "N" is passed by value (both ways - i.e. into the called procedure and also copied by value back into the calling routine).
- A6 is used as the frame pointer (assumed to have initial value \$00002000).
- Procedure "Process" uses registers D0 - D4 as well as registers A0 - A4.
- Array X starts at location \$1800, Y starts at \$17F8, Z is at \$17FC, ALPHA is at \$17FD, and N is at \$17FE.

Recursive Subroutine Calls Example

Problem specification (continued):

{main routine}

X: array [0..30] of words

Y: longword

Z, ALPHA, N: byte

Process(var: X, var: Y, var: Z, var: ALPHA, N)

- **We are to show all the M68000 assembly language instructions necessary to pass these parameters as well as to copy the return value N into its regular storage location (off the stack) (at \$17FE).**

Recursive Subroutine Calls Example

Problem specification (continued):

Procedure Process (A, B, C, D, E)

A: array [0..?] of words {passed by reference}

B: longword {passed by reference}

C, D: byte {passed by reference}

E: byte {passed both ways by value}

local variables -

T: longword

U: word

V: byte

{ some place within the first invocation of "Process" it calls itself as follows: }

Process(var: A, var: T, var: C, var: V, E) {Note that some input parameters are passed through to the next iteration.}

Recursive Subroutine Calls Example

Solution

The main program is assumed to allocate the original storage for:

	ORG \$17F8	
Y	DS.L 1	This will resolve to address \$000017F8
Z	DS.B 1	This will resolve to address \$000017FC
ALPHA	DS.B 1	This will resolve to address \$000017FD
N	DS.B 1	This will resolve to address \$000017FE
*		
	ORG \$1800	
X	DS.W 31	an array of longwords 0..30

Recursive Subroutine Calls Example

Solution (Continued)

ORG \$1000 (assumed where main program starts - not critical)

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* In main program the procedure (subroutine) is called in HLL:

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* Process (var:X, var:Y, var:Z, var:ALPHA, N) where N is the only one passed by value

* The assembly language version/translation of this invocation is:

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CLR.W D2	zeroes out an entire word for pushing on stack
MOVE.B N,D2	copies value of byte N into lowest byte of D2
MOVE.W D2,-(A7)	pushes that word containing value of N on stack
PEA ALPHA	pushes pointers to other arguments in reverse
PEA Z	order
PEA Y	
PEA X	
JSR Process	actually call the subroutine here
MOVE.B 17(A7),N	copy returned value back into N
ADDA.L #18,A7	fix up stack from all parameters pushed for
	subroutine call.

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Recursive Subroutine Calls Example

Solution (Continued) Stack Utilization Diagram

0E5E	not used	0E94	local 2 "T"	0ECA	A0
0E60			(longword)		
0E64	not used	0E98	local 2 "U"	0ECE	A1 (high)
		0E9A	- - "v" 2		A1 (low)
0E68	not used	** 0E9C	link reg val	0ED2	A2
			= \$00000EE6		
0E6C	D0 (high) 2	0EA0	return addr	0ED6	A3
	D0 (low)		into Process		
0E70	D1 2	0EA4	Addr of "X"	0EDA	A4
			= "A" in Proc		
0E74	D2 2	0EA8	Addr of "T"1	0EDE	local 1 "T"
			= \$00000EDE		(longword)
0E78	D3 2	0EAC	Addr of "Z"	0EE2	local 1 "U"
			equiv "C" 1	0EE4	- - "v" 1
0E7C	D4 2	0EB0	Addr of "V"1	*0EE6	orig linkreg
			= \$00000EE5		= \$00002000
0E80	A0 2	0EB4	\$00 "E"2	0EEA	return addr
		0EB6	D0 (high) 1		into main pr
0E84	A1 2		D0 (low)	0EEE	Addr of "X"
		0EBA	D1 1		= \$00001800
0E88	A2 2	0EBE	D2 1	0EF2	Addr of "Y"
					= \$000017F8
0E8C	A3 2	0EC2	D3 1	0EF6	Addr of "Z"
					= \$000017FC
0E90	A4 2	0EC6	D4 1	0EFA	Addr "ALPHA"
					= \$000017FD
				0EFE	\$00 "N"val

* indicates the value of link register A6 during first call of Process

** indicates the value of link register A6 during the second call to Process

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Recursive Subroutine Calls Example Solution

(Continued) procedure Process

- The coding of procedure Process would be something like this:

Procedure Process (var:A, var:B, var:C, var:D, E)

* where A: is an array of words [0.. ?] passed by reference

* B: longword passed by reference

* C, D: byte passed by reference

* E: byte passed by value (in BOTH directions)

* and local variables:

* T: longword

* U: word

* V: byte

Aptr	equ	8	displacements for finding pass by reference
Bptr	equ	12	addresses from the frame pointer: A6
Cptr	equ	16	
Dptr	equ	20	
E	equ	25	this one is a byte which is passed by value
V	equ	-1	
U	equ	-4	
T	equ	-8	

Recursive Subroutine Calls Example

Solution (Continued) procedure Process

* The start of the code of Process looks like this:

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```
Process  LINK    A6,#-8
         MOVEM.L D0-D4/A0-A4,-(A7)      save registers as required
```

*

* The invocation of Process from within Process:

*

* Process (A, T, C, V, E)

*

```
        CLR.W    D0
        MOVE.B   E(A6),D0              note how we access "E" - we could have
        MOVE.W   D0,-(A7)             modified "E" before sending it
        PEA      V(A6)                this is basically how we can use "V" too
        MOVE.L   Cptr(A6),-(A7)       we push the pointer to "Z" on stack
        PEA      T(A6),A0             push pointer to local variable "T" on stack
        MOVE.L   Aprt(A6),-(A7)       push pointer to "X" ("A" in Process)
        JSR      Process
        MOVE.B   17(A7),E(A6)         copy return value of "E" to local copy
        ADDA.L   #18,A7               fix up stack from all parameters pushed
```

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Recursive Subroutine Calls Example

Solution (Continued) procedure Process

* This is how we'd access some of the variables in the subroutine:

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MOVEA.L Aptr(A6),A0	This is how we'd copy the first array
MOVE.L (A0),U(A6)	element of X ("A" in procedure) into "U"

*

MOVEA.L Bptr(A6),A1	This is how we'd copy input parameter "B"
MOVE.W (A1),T(A6)	into local word "T"

*

MOVEA.L Cptr(A6),A2	This is how we actually reference "C"
MOVE.B (A2),D1	

*

MOVEA.L Dptr(A6),A3	This is how we could access/change
CLR.B (A3)	"D" in procedure = "ALPHA" in main

*

* Before leaving the procedure we'd need to restore registers and destroy stack frame:

*

```
MOVEM.L (A7)+,D0-D4/A0-A4  
UNLK A6  
RTS
```